

1. Each item is 2.5 points.

- (a) Quotient 5, remainder $-1 - 2i$.
 - (b) Need $a^2 + 2b^2 = 9$ yielding ± 3 and $\pm 1 \pm 2\sqrt{-2}$.
 - (c) Take the smallest power of each irreducible factor: gcd is $3(2 - i)$.
 - (d) The classes are represented by polynomials of degree ≤ 2 , so there are 7^3 residue classes.
 - (e) Units are $\bar{1}, \bar{2}, \overline{x+1}, \overline{2x+2}$; zero divisors are $\bar{x}, \overline{x+2}, \overline{2x}, \overline{2x+1}$.
 - (f) Total is $\frac{1}{7}(2^7 - 2) = 18$.
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2. Each item is 4 points.

- (a) Solution is $p \equiv x + 2x^2 \pmod{x^3 - 2x^2}$.
 - (b) There are primitive roots mod 34 and 37 but not mod 35 or mod 36.
 - (c) 2 is a primi root mod 3^{2026} so $2+3^{2026}$ is a prim root mod $2 \cdot 3^{2026}$. Total number is $\varphi(\varphi(2 \cdot 3^{2026})) = 2 \cdot 3^{2026}$.
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3. Part (a) is 2 points, (b) is 3 points, (c) and (d) are 4 points each.

- (a) The residue classes are represented by polynomials of degree less than 3: $\bar{0}, \bar{1}, \bar{x}, \overline{x+1}, \overline{x^2}, \overline{x^2+1}, \overline{x^2+x}, \overline{x^2+x+1}$.
 - (b) We have $\overline{x^2+x^2+1} = \bar{1}$, $\overline{x^2 \cdot x^2+1} = \overline{x^2+1}$, and $\overline{x^2+1}^2 = \bar{0}$.
 - (c) The units are the polynomials relatively prime to the modulus: $\bar{1}, \bar{x}, \overline{x^2}, \overline{x^2+x+1}$.
The zero divisors are the nonzero polynomials not relatively prime to the modulus: $\overline{x+1}, \overline{x^2+1}, \overline{x^2+x}$.
 - (d) There are 4 units and indeed $\overline{x^2+x+1}^4 = \overline{x^2} = \bar{1}$ as required.
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4. Part (a) is 2 points, and (b) and (c) are 5 points each.

- (a) Note $N(2 + \sqrt{-7}) = 11$ is a prime integer, so $2 + \sqrt{-7}$ is irreducible.
 - (b) Note $N(1 + \sqrt{-7}) = 8$ so if we had a nontrivial factorization, it would have to be the product of an element of norm 2 with an element of norm 4. But since $N(a + b\sqrt{-7}) = a^2 + 7b^2$ there are no elements of norm 2 or 4, so there is no possible factorization.
 - (c) Note that $(1 + \sqrt{-7})(1 - \sqrt{-7}) = 8 = 2 \cdot 4$ so $1 + \sqrt{-7}$ divides $2 \cdot 4$ but it divides neither 2 nor 4, since $2/(1 + \sqrt{-7}) = (1 - \sqrt{-7})/4$ and $4/(1 + \sqrt{-7}) = (1 - \sqrt{-7})/2$. This means $1 + \sqrt{-7}$ is not prime.
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5. For $p(x) = x^3 + x + 1$ we have $p(0) = p(2) = p(3) = 1$, $p(1) = 3$, $p(4) = 4 \pmod{5}$, so p has no roots. Since it has degree 3 it is irreducible, so $\mathbb{F}_5[x]$ modulo $x^3 + x + 1$ is a field.

6. 3 is a primitive root modulo 7 since its order divides $\varphi(7) = 6$ but $3^2, 3^3 \not\equiv 1 \pmod{7}$ so its order is 6. We then compute $3^6 \equiv 43 \pmod{49}$: thus 3 is a primitive root mod 7^2 hence mod 7^d for all $d \geq 2$.

- 7. (a) Using a zero-knowledge protocol like the Rabin protocol described in class, where Peggy proves to an arbitrarily high probability that she knows the square root of a particular value s^2 modulo $N = pq$, will allow Peggy to convince Victor that she knows the secret s without revealing useful information.
 - (b) Using primality/compositeness tests like the Fermat test, the Lucas primality criterion, Miller-Rabin, or AKS allow for rapid and accurate testing of primality even for very large integers.
 - (c) Trial division will be very slow for 30-digit integers, but Pollard ρ and the sieving methods will allow us to factors of that size quite rapidly (Pollard ρ typically takes $\sim N^{1/4}$ time to factor N , which for $N \approx 10^{30}$ gives a computation size of $\approx 10^{7.5}$ steps, very doable).
 - (d) If a reducible polynomial $p(x)$ has degree 3, it must factor as a product of three degree-1 terms, or a degree-2 term and a degree-1 term. But any degree-1 term $ax - b$ yields a root $x = b/a$ of $p(x)$. So, if $p(x)$ has no roots, it has no possible factorization.
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1. Each part is 6 points. Use the Euclidean algorithm and then solve for each of the remainders.

- (a) GCD is $x^2 + x$, with linear combination $1 \cdot (x^4 + x) + x \cdot (x^3 + x) = x^2 + x$.
 - (b) GCD is 1, with linear combination $(1 - 2i)(9 - 5i) + (4 + 5i)(3 + 2i) = 1$.
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2. Each part is 3 points.

- (a) Inverse of $1 + i$ is $-4 + 3i$ (Euclid) so solution is $n \equiv 3(-4 + 3i) \pmod{8 + i}$.
 - (b) Searching for roots produces factorizations $x(x + 1)$, $(x + 1)^2$, and $(x + 2)^2$.
 - (c) 2 is a primitive root mod 3^2 hence mod 3^{2026} . Total number is $\varphi(\varphi(3^{2026})) = 2 \cdot 3^{2026}$.
 - (d) Total is $\frac{1}{4}(7^4 - 7^2) = 588$.
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3. Each part is 4 points.

- (a) Note $N(7 + 4\sqrt{3}) = 1$ so it is a unit since the norm is ± 1 . The inverse is the conjugate $7 - 4\sqrt{3}$.
 - (b) Note $N(4 + 5i) = 4^2 + 5^2 = 41$ is a prime integer so as $\mathbb{Z}[i]$ is Euclidean, $4 + 5i$ is irreducible and prime.
 - (c) Searching yields a root $x = 3$, so the polynomial is reducible so $\mathbb{F}_5[x]$ modulo $x^4 + x + 1$ is not a field.
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4. Since $3 + 2i$ is irreducible (as it has norm 13, which is prime) and there are 13 residue classes, we see that $\mathbb{Z}[i]$ modulo $3 + 2i$ is a field. Then we see $i^{13} \equiv i \pmod{3 + 2i}$ as required (indeed, i^{13} just equals i).

5. $4 = 2 \cdot 2 = (1 + \sqrt{-3})(1 - \sqrt{-3})$, not equivalent since only units are ± 1 . As $N(2) = N(1 \pm \sqrt{-3}) = 4$ and $N(a + b\sqrt{-3}) = a^2 + 3b^2$ there are no elements of norm 2, so any element of norm 4 is irreducible. Euclidean domains have unique factorization so $\mathbb{Z}[\sqrt{-3}]$ cannot be Euclidean.

6. It is not hard to list all the units to see that there are 4 of them (they are the polynomials with constant term 1). We then calculate $x^2 + 1^4 = \frac{x^4 + 2x^2 + 1^2}{x^4 + 2x^2 + 1^2} = \bar{1}^2 = \bar{1}$ so Euler's theorem holds.

7. Since $125 = 5^3$ we can use $\mathbb{F}_5[x]$ modulo an irreducible polynomial of degree 3. We actually just identified such a polynomial, namely $x^3 + x + 1$, in part (j).

8. Each part is 3 points.

- (a) Using the Fermat factorization method allows us to factor an integer $N = pq$ where p, q are close together by testing whether $a^2 - N$ is a square for $a > \sqrt{N}$.
 - (b) Among the various factorization algorithms discussed in class like trial division, Pollard $p - 1$, Pollard ρ , and the sieving methods, none allows for extremely fast factorization of large integers (factoring integers more than 100 base-10 digits takes a huge amount of time and memory).
 - (c) Because the ring $\mathbb{Z}[i]$ of Gaussian integers is a Euclidean domain, as proven in class, the Euclidean algorithm can be used to compute greatest common divisors. The norm of the remainder at each step is at most half the norm of the value being divided by, so the algorithm is very efficient.
 - (d) Polynomials of degree greater than 3 could have a factorization where each of the irreducible factors has degree greater than 1. For instance, $x^4 + 3x^2 + 2$ factors over the real numbers as $(x^2 + 1)(x^2 + 2)$, and both factors are irreducible since they have no real roots.
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