Math 4527 (Number Theory 2) Lecture #1 of 38 \sim January 20, 2021

Introduction + Pythagorean Triples

- Pythagorean Triples (part 2)
- Linear Diophantine Equations
- The Frobenius Coin Problem

This material represents $\S6.1.1-6.1.3$ from the course notes.

More Pythagorean Triples, I

Recall from last time:

Definition

We say a Pythagorean triple (x, y, z) satisfying $x^2 + y^2 = z^2$ is <u>primitive</u> if gcd(x, y, z) = 1.

Theorem (Primitive Pythagorean Triples)

Every primitive Pythagorean triple of the form (x, y, z) with x even is of the form $(x, y, z) = (2st, s^2 - t^2, s^2 + t^2)$, for some relatively prime integers s > t of opposite parity. Conversely, any such triple is Pythagorean and primitive.

I gave two proofs last time (one using arithmetic in \mathbb{Z} and another using arithmetic in $\mathbb{Z}[i]$). Now I will do the third proof, which uses geometry.

More Pythagorean Triples, II

<u>Proof #3</u>:

- Suppose $x^2 + y^2 = z^2$ and x, y, z are relatively prime.
- Dividing by z^2 yields the equivalent equation $\left(\frac{x}{z}\right)^2 + \left(\frac{y}{z}\right)^2 = 1.$
- Each Pythagorean triple (x, y, z) therefore yields a point (a, b) = (x/z, y/z) with rational coordinates on the unit circle x² + y² = 1.
- Conversely, if we have a point (a, b) with rational coordinates on the unit circle, then if z is the lcm of the denominators so that (a, b) = (x/z, y/z), we obtain a primitive triple (x, y, z)with $x^2 + y^2 = z^2$.
- Therefore, finding the primitive Pythagorean triples is equivalent to describing all points (a, b) on the unit circle x² + y² = 1 whose coordinates are both rational numbers.

More Pythagorean Triples, III

<u>Proof #3</u> (continued):

• To do this, consider a line passing through the point (-1, 0) with a finite slope:



- Such a line will intersect the circle x² + y² = 1 in exactly one other point.
- If the coordinates of this point are rational, then the line will have rational slope.
- Conversely, if the line has rational slope, its other intersection point with the circle will be rational.

<u>Proof #3</u> (continued more):

- Explicitly, if the line has slope $\frac{t}{s}$, its equation is $y = \frac{t}{s}(x+1)$.
- We can then simply plug in $y = \frac{t}{s}(x+1)$ to $x^2 + y^2 = 1$ and solve to see that the other intersection point is $(x, y) = \left(\frac{s^2 t^2}{s^2 + t^2}, \frac{2st}{s^2 + t^2}\right)$, which is rational.
- Thus, the rational points on the unit circle are those of the form $\left(\frac{s^2 t^2}{s^2 + t^2}, \frac{2st}{s^2 + t^2}\right)$ for some integers s and t.
- Clearing the denominator immediately yields the desired Pythagorean triples.

If you've never seen it before, the idea in the third proof actually has a surprisingly useful application in calculus.

- Specifically, if we write u = s/t, then we obtain a rational parametrization of the unit circle in terms of u.
- Explicitly, we have $\cos \theta = \frac{1-u^2}{1+u^2}$ and $\sin \theta = \frac{2u}{1+u^2}$.
- Using some elementary geometry, it is not hard to see that this parameter is $u = \tan(\theta/2)$.
- The upshot: if we make the substitution $u = \tan(\theta/2)$, then both $\sin \theta$ and $\cos \theta$ are rational functions in u.
- This substitution is called the <u>Weierstrass substitution</u> and can be used to evaluate integrals like $\int \frac{1}{3+\sin\theta} d\theta$ that are quite hard otherwise.

Using the characterization above, we can easily generate a list of Pythagorean triples with hypotenuse \leq 80:

5	t	Primitive Triple	Other Triples		
2	1	(3,4,5)	(6,8,10), (9,12,15),, (48,64,80)		
3	2	(5, 12, 13)	(10, 24, 26),, (30, 72, 78)		
4	1	(8, 15, 17)	(16, 30, 34),, (32, 60, 68)		
4	3	(7, 24, 25)	(14, 48, 50), (21, 72, 75)		
5	2	(20, 21, 29)	(40, 41, 58)		
5	4	(9, 40, 41)			
6	1	(12, 35, 37)	(24, 70, 74)		
6	5	(11, 60, 61)			
7	2	(28, 45, 53)			
7	4	(33, 56, 65)			
8	1	(16, 63, 65)			
8	3	(48, 55, 73)			

<u>Example</u>: Find all Pythagorean right triangles having one side of length 20.

• Any such right triangle has legs of lengths k(2st) and $k(s^2 - t^2)$, with hypotenuse $k(s^2 + t^2)$, where s > t are unique positive integers of opposite parity and k is some unique positive integer. Now we just try each possibility:

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- 1. Suppose 20 = 2stk.
 - Then 10 = stk, so (s, t, k) = (10, 1, 1) or (5, 2, 1).
 - This yields 20-99-101 and 20-21-29 triangles.

Example (continued): Find all Pythagorean right triangles having one side of length 20.

- 2. Suppose $20 = k(s^2 t^2)$.
 - Then k must be divisible by 4 since $s^2 t^2$ is odd.
 - Since $k \neq 20$ we see k = 4. Then $s^2 t^2 = 5$ requires s = 3 and t = 2.
 - This yields a 15-20-25 triangle.

3. Suppose $20 = k(s^2 + t^2)$.

- Then since s² + t² ≥ 5, the only possibilities are k = 4 (then s = 2, t = 1), k = 2 (then s = 3, t = 1 but these are not of opposite parity) or k = 1 (then s = 4 and t = 2 but again these don't work).
- This yields a 12-16-20 triangle.

So we get 4 triples: (20,99,101), (20,21,29), (15,20,25), (12,16,20).

We will now take a step in the "easier" direction and talk about the simpler class of linear Diophantine equations.

- Solving a linear equation in one variable over the integers is trivial (the solution to ax = b is x = b/a, assuming a is nonzero and divides b).
- So the simplest interesting equations are linear equations in two variables.
- The general form of a linear equation in two variables is ax + by = c, for some fixed integers *a*, *b*, and *c*.
- Our goal is to determine when this equation has an integral solution (x, y), and then to characterize all the solutions.
- In keeping with our theme of "exploiting other rings to solve problems in \mathbb{Z} ", we will use modular arithmetic to reduce the two-variable equation to a one-variable equation.

Linear Diophantine Equations, II

We will use the following proposition about linear congruences modulo m:

Proposition (Linear Equations Mod m)

The equation $ax \equiv b \pmod{m}$ has a solution for x if and only if d = gcd(a, m) divides b.

If d|b, then the set of all such x is given by the residue class \overline{r} modulo m/d, where r is any solution to the equation.

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Examples:

- 1. The equation $9x \equiv 5 \mod 12$ has no solutions, because gcd(9, 12) = 3 does not divide 5. The point is: any multiple of 9 will always be divisible by 3 modulo 12, so it can't equal 5.
- 2. The equation $9x \equiv 6 \mod 12$ has solutions, since gcd(9, 12) = 3 does divide 6. Since x = 2 is a solution, the full set of solutions is $x \equiv 2 \mod 4$.

Linear Diophantine Equations, III

Proof:

- If x is a solution to the congruence ax ≡ b (mod m), then there exists an integer k with ax - mk = b. Since d = gcd(a, m) divides the left-hand side, it must divide b.
- Now suppose d = gcd(a, m) divides b, and set a' = a/d, b' = b/d, and m' = m/d.
- Then the original equation becomes a'dx ≡ b'd (mod m'd), which is equivalent to a'x ≡ b' (mod m'). This is a property of congruences that is easy to verify if you write it in terms of divisibility.
- But since a' and m' are relatively prime, a' is a unit modulo m', so we can simply multiply by its inverse to obtain x ≡ b' ⋅ (a')⁻¹ (mod m'). This means that there is a unique solution to the congruence modulo m' = m/d, as claimed.

Linear Diophantine Equations, IV

By reducing modulo one of the coefficients, we can solve linear Diophantine equations in two variables:

Theorem (Linear Diophantine Equations in 2 Variables)

Let a, b, c be integers with $ab \neq 0$, and set d = gcd(a, b). If $d \nmid c$, the equation ax + by = c has no integer solutions (x, y). If $d \mid c$, ax + by = c has infinitely many integer solutions (x, y). If (x_0, y_0) is one solution, then all the others are $(x_0 - bt/d, y_0 + at/d)$, for some integer t.

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Examples:

- 1. The equation 9x + 12y = 5 has no integer solutions, since gcd(6,9) = 3 does not divide 5.
- The equation 9x + 12y = 6 has integer solutions because gcd(9, 12) = 3 does divide 6. One solution is (x, y) = (2, -1), so the full set is (x, y) = (2 4t, -1 + 3t) for integers t.

Proof:

- If a = b = 0, then the equation ax + by = c is either trivially true (if c = 0) or trivially false (if $c \neq 0$), so we can assume that the gcd d is nonzero.
- If one of a, b is zero, the equation is also trivial, so we may also deal only with the case where $ab \neq 0$.
- In this case, observe that there is an integral solution to ax + by = c if and only if there is a solution to the congruence $ax \equiv c \pmod{b}$, since then $y = \frac{c ax}{b}$.
- From our proposition above, we know that ax ≡ c (mod b) has a solution only if d = gcd(a, b) divides c. This is the first part of the theorem.

<u>Proof</u> (continued):

- For the second part, suppose d = gcd(a, b) does divide c, and consider the values of x satisfying ax + by = c.
- In this case, again by the proposition we just proved, if we set a' = a/d, b' = b/d, and c' = c/d, the set of all such x is given by the residue class x
 ₀ modulo b', where x₀ ≡ c' ⋅ (a')⁻¹ (mod b').
- Now if (x, y) is any solution, then by the above, we see that x = x₀ bt/d for some integer t, and then we can directly compute y = y₀ + at/d where ax₀ + by₀ = c.
- Since these are all solutions, this yields the full characterization of the solutions given above.

Example: Find all solutions to 14x + 18y = 12 in integers (x, y).

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- First, we compute gcd(14, 18) = 2, and then divide through by the gcd to get 7x + 9y = 6.
- This is equivalent to solving $7x \equiv 6 \pmod{9}$.
- We compute (via the Euclidean algorithm or guess-and-check) that the inverse of 7 mod 9 is 4, so multiplying both sides by 4 yields x = 24 = 6 (mod 9).
- Hence one solution is (x, y) = (6, -4). The set of all solutions is then $(x, y) = \boxed{(6 9t, -4 + 7t)}$ for $t \in \mathbb{Z}$.

<u>Example</u>: Find all solutions to 372x + 948y = 42 in integers (x, y).

- Using the Euclidean algorithm we can quickly compute gcd(372,948) = 12.
- Since 12 does not divide 42, there are no solutions.

We can solve linear Diophantine equations in two variables by making changes of variable. Here's an example:

<u>Example</u>: Find all solutions to the equation 4x + 13y = 5.

- By the division algorithm, we have 13 = 3 · 4 + 1, so we can write the system in the form 4x + (3 · 4 + 1)y = 5, and rearrange this into the form 4(x + 3y) + 1y = 5.
- If we substitute u = x + 3y, this new system becomes 4u + y = 5, which we can easily solve to get y = 5 4u.
- Substituting back yields x = u 3y = u 3(5 4u) = -15 + 13u.
- Thus, we obtain the general solution (x, y) = (-15 + 13u, 5 4u).

This method, using changes of variable, is the most efficient way to solve systems of linear Diophantine equations involving more variables or equations.

- The approach is essentially the same as the standard linear algebra procedure of row-reducing a matrix to solve a system of equations.
- The easiest approach is convert the system into matrix form, and then perform row and column operations on the matrix until it is in a sufficiently simple form that the solution to the original system is obvious.
- The general procedure for solving a system of linear equations over Z is essentially the same, except for the added complication that all of the row and column operations need to be done over Z.

As with a system of equations over a field, the end result will be either that the system has no solution, a unique solution, or an infinite family of solutions with some number of free parameters.

- I won't bother going into the technical details or proving that this procedure always works, since it really is a topic from abstract algebra.
- For those curious: it is in fact equivalent to the procedure for converting a presentation of a finitely generated additive abelian group into a description of the abelian group as a direct product of cyclic groups, which is in turn a special case of the general classification theorem for finitely-generated modules over a principal ideal domain.

So I'll just give an example and let you see the general idea.

Example: Find all solutions to 3x + 7y + 8z = 13 in integers (x, y, z).

- Motivated by the division algorithm, we rewrite the equation as 3(x + 2y + 2z) + y + 2z = 13, and then substitute w = x + 2y + 2z.
- The new equation is 3w + y + 2z = 13, which we can easily solve for y, yielding y = 13 3w 2z.
- Then x = w 2y 2z = 7w + 2z 26.
- So we obtain the general solution

 (x, y, z) = (7w + 2z 26, 13 3w 2z, z) where w and z
 are arbitrary integers.

In various settings (some of which are actually motivated by real-world concerns for once!), we are sometimes also interested in knowing for which values of c the equation ax + by = c has a solution in *nonnegative* integers (x, y).

- For example, if there are postage stamps worth 5 cents and stamps worth 13 cents, is it possible to use them to put exactly 79 cents' worth of postage on an envelope? (Here we want to solve 5x + 13y = 79.)
- Another version occurs in sports: In American football, a team can score 3 points for a field goal, or 7 points for a touchdown. What possible scores can a team obtain? (Ignore safeties, missed extra points, and so forth.)

We remark that we can reduce to the situation with a, b relatively prime by dividing through by their gcd.

The most obvious method is simply to make a list of totals that are attainable.

- For example, for 5x + 13y, we obtain values 0, 5, 10, 13, 15, 18, 20, 23, 25, 26, 28, 30, 31, 33, 35, 36, 38, 39, 40, 41, 43, 44, 45, 46, 48, 49, 50, 51, 52, 53, 54, 55, 56, 57, 58,
- Likewise, for 3x + 7y we obtain values 0, 3, 6, 7, 9, 10, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21, 22, 23,
- In each case, it seems like we miss only finitely many values.

The problem of describing the largest integer that cannot be written as a nonnegative linear combination of two integers (also called the <u>Frobenius coin problem</u>) was first solved by Sylvester:

Theorem (Sylvester)

If a and b are relatively prime integers, then there are exactly $\frac{1}{2}(a-1)(b-1)$ integers that cannot be written in the form ax + by with $x, y \ge 0$, and the largest such integer is ab - a - b.

<u>Remark</u>: In mathematics competition circles, this result is often known as the "Chicken McNuggets Theorem".

The Frobenius Coin Problem, IV

Proof:

- For brevity, we say an integer is "representable" if it can be written in the form ax + by with x, y ≥ 0.
- Without loss of generality, assume *a* < *b*. Arrange the nonnegative integers in an array in the following manner:

0	1	2	•••	a — 1
а	a+1	a + 2	• • •	2 <i>a</i> — 1
2 <i>a</i>	2a + 1	2 <i>a</i> + 2	•••	3 <i>a</i> — 1
÷	:	:		÷
ab — a	ab-a+1	<i>ab</i> — <i>a</i> + 2		<i>ab</i> — 1

• Now we use the array to mark all of the representable integers. We first box all of the multiples of *b*: then an integer is representable precisely if it appears in the same column as some multiple of *b*, lower down.

Proof (continued):

• For illustration, here is the array with a = 4 and b = 7:



Proof (continueder):

- Since a and b are relatively prime, the integers 0, b, 2b, ..., (a-1)b all lie in different columns. Thus, the largest element that is left unmarked is the element one row above (a 1)b, which is ab a b, so this is the largest integer not expressible as ax + by with x, y ≥ 0.
- For the other part, we simply count the number of unmarked integers in the array.
- The number of integers lying above kb is $\lfloor kb/a \rfloor$, so there are

a total of
$$\sum_{k=0}^{a-1} \left\lfloor \frac{kb}{a} \right\rfloor$$
 unmarked integers in the array.

The Frobenius Coin Problem, VII

<u>Proof</u> (continuederer):

- We can interpret the sum $\sum_{k=0}^{a-1} \lfloor \frac{kb}{a} \rfloor$ geometrically as the number of lattice points lying under the line y = (b/a)x, with $1 \le x \le a 1$.
- Equivalently, this is the total number of lattice points lying strictly inside the rectangle with vertices (0,0), (a,0), (a,b), (0,b) and below the diagonal.
- By symmetry, since there are no lattice points on the interior of the diagonal, exactly half of the lattice points inside the a × b rectangle are below the diagonal.
- Since this full set of points forms an $(a-1) \times (b-1)$ rectangle, there are (a-1)(b-1) such lattice points. Therefore, the number of unmarked integers in the array is $\frac{1}{2}(a-1)(b-1)$, as claimed.

<u>Example</u>: There are postage stamps worth 5 cents and stamps worth 13 cents. What is the largest non-attainable amount of postage, and how many non-attainable amounts are there?

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- By Sylvester's theorem with a = 5 and b = 13, the largest non-representable integer is $5 \cdot 13 13 5 = 47$.
- In total, there are $\frac{1}{2} \cdot 4 \cdot 12 = 24$ unattainable totals.

We could of course generalize this problem, to ask: for given integers $a_1, a_2, ..., a_k$, what is the largest integer *n* that cannot be written as a nonnegative integer linear combination of the a_i ?

- It turns out that there is no known general formula when k > 2 (though the result is fairly effectively computable for k = 3).
- For a fixed number of denominations k, there does exist a polynomial-time algorithm (polynomial in log a_k, specifically) for computing this maximum integer n, but it is not appreciably faster than merely attempting to list the possibilities!
- For a variable number of denominations *k*, it is known that computing *n* is *NP*-hard.



We discussed a bit more about Pythagorean triples.

We discussed how to solve linear Diophantine equations in two variables and made some remarks about larger systems.

We discussed the Frobenius coin problem.

Next lecture: The Farey sequences and rational approximation.